

Optimizing Age of Information with Attacks

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- 1 Motivation
- 2 Model Description
- 3 Main Results
 - Non-preemptive M/M/1/1 Queue
 - Preemptive M/M/1/1 Queue
 - Preemptive M/M/1/1 Tandem Queues
 - Non-preemptive M/M/1/1 Tandem Queues
- 4 Conclusions

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Introduction

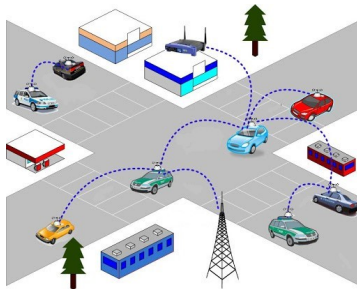
In a system that is observed remotely (by a monitor)

Timely information is crucial

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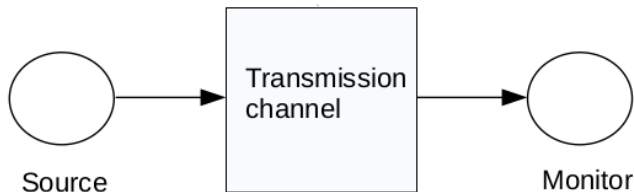
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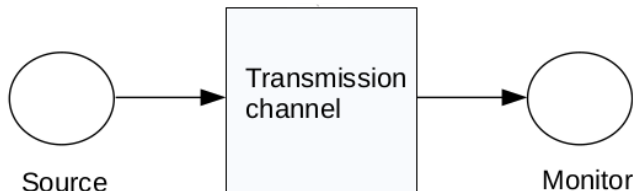
Other examples in industry

Age of Information



Generated packets (updates): information + timestamp

Age of Information



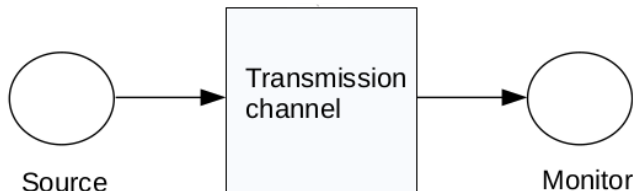
Generated packets (updates): information + timestamp

Definition: Age of Information (AoI)

The time elapsed since the generation time of the last arrived packet at the destination.

Large AoI \Rightarrow obsolete information

Age of Information



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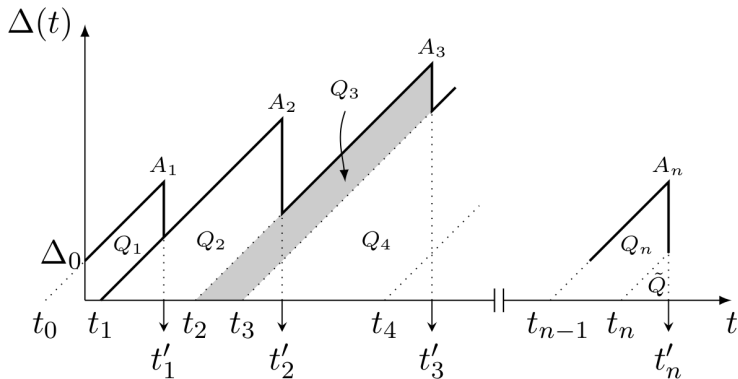
Large AoI \Rightarrow obsolete information

Small AoI \Rightarrow **timely (or recent)** information!

Example

t_i : generation time of packet i

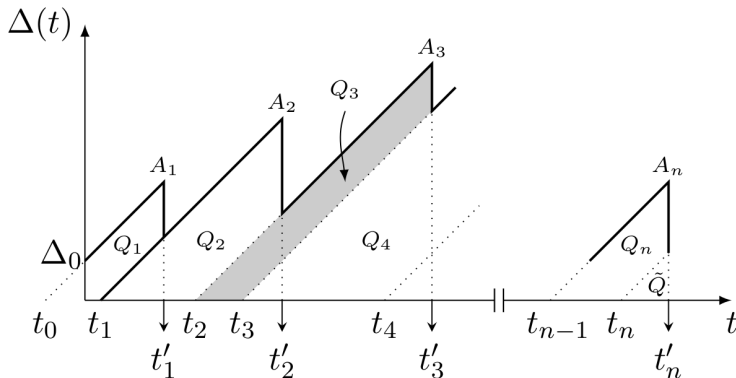
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Average Aol calculation:

Example

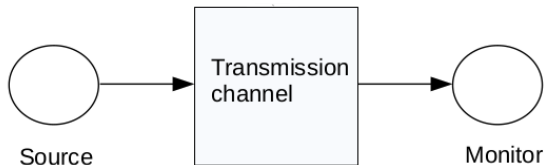
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Average Aol calculation: not possible in a general setting

Age of Information

Average AoI calculation: not possible in a general setting

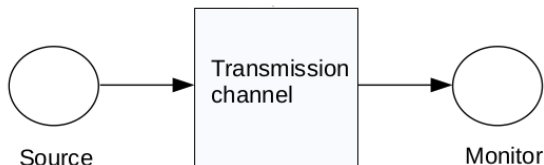


For particular queueing models (of the transmission channel)

Known expressions for M/M/1, M/M/1/1, M/M/1/2, M/M/2...

Age of Information

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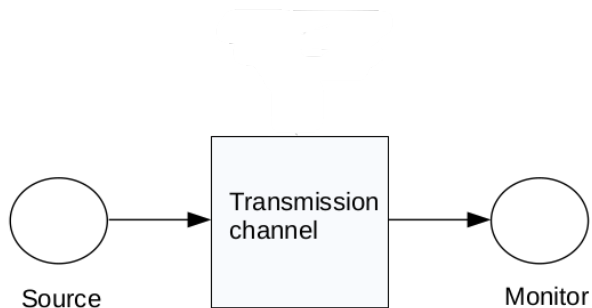


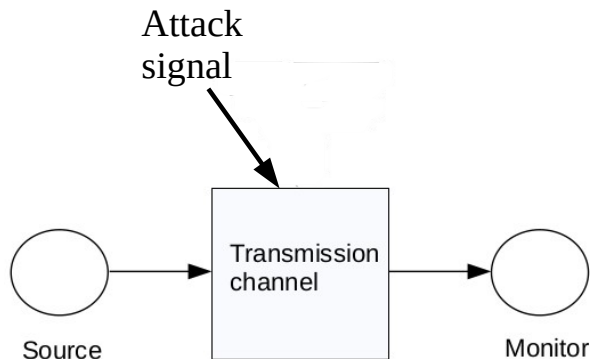
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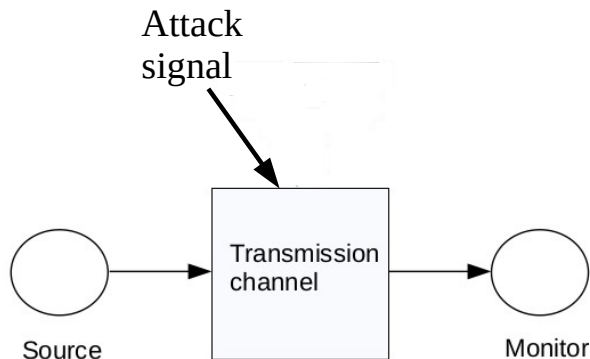
Current challenges

- Analytical expression of interesting queueing models
- Optimization problems with derived expressions





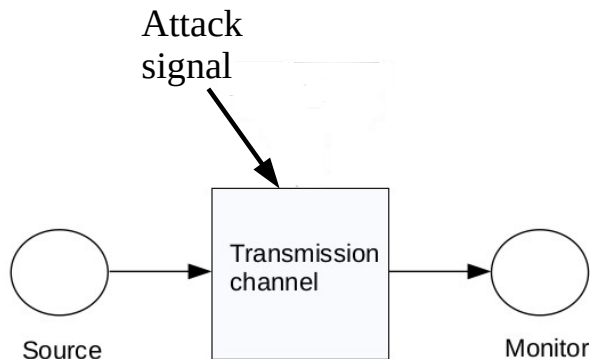
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Questions:

Do attacks minimize Aol?



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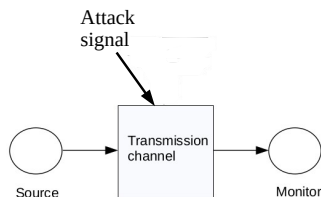
Do attacks minimize Aol?

Why and under which conditions?

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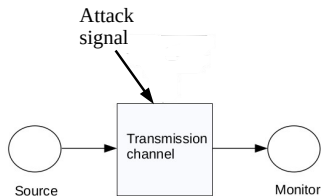
Model Description

Generation time of updates: Poisson with rate λ
Service time in the queues: exponential with rate μ
Attacks signal: Poisson with rate α



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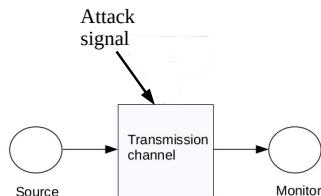
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$\Delta(\alpha)$: Average Aol as a function of α

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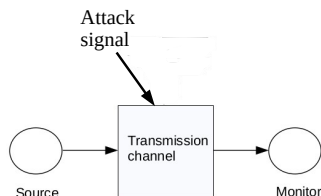


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$$\mathcal{R} = \frac{\lim_{\alpha \rightarrow 0+} \Delta(\alpha)}{\min_{\alpha} \Delta(\alpha)}$$

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$\mathcal{R} > 1 \Rightarrow$ Attacks minimize Aol

Considered Models

	M/M/1/1 Queue	Tandem M/M/1/1 Queues
Non-preemptive		
Preemptive		

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Non-preemptive M/M/1/1 Queue

Proposition

$$\Delta(\alpha) = \frac{1}{\lambda} + \frac{1}{\mu} + \frac{\lambda}{(\mu + \alpha)(\lambda + \mu + \alpha)} + \frac{\alpha}{\lambda\mu}.$$

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$\mathcal{R} > 1$ if and only if $\lambda/\mu > 1,2469$

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- Preemptive M/M/1/1 Queue: attacks DO NOT reduce AoI

Preemptive M/M/1/1 Tandem Queues

n : number of preemptive M/M/1/1 queues in tandem
Service times in queue i : $\text{Exp}(\mu_i)$

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where $\mu_0 = \lambda$.

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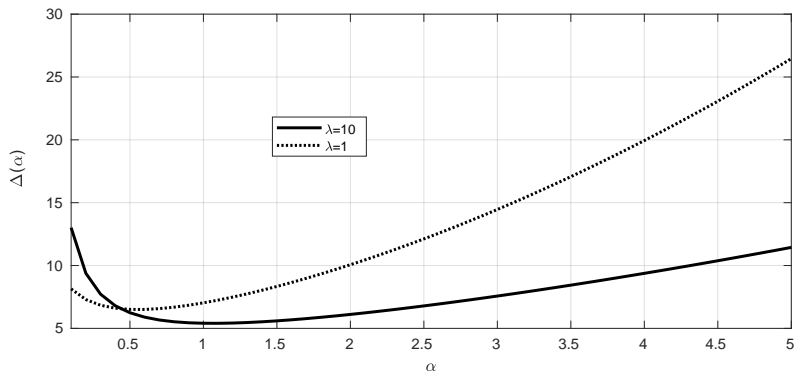
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- Preemptive M/M/1/1 Tandem Queues DO NOT reduce AoI

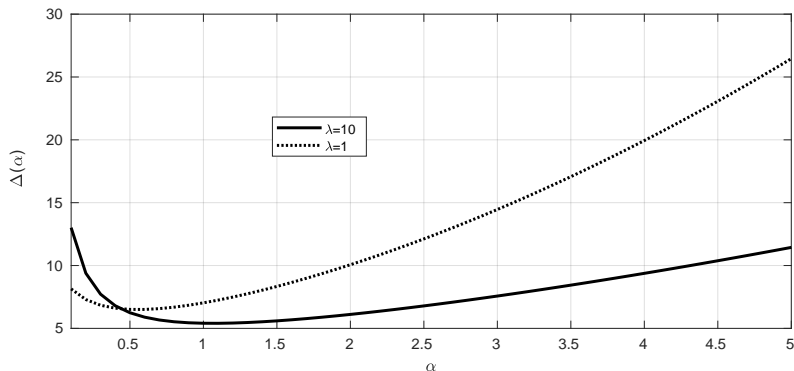
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$$\mu_1 = \mu_2 = 1$$



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$\mathcal{R} = 1$ for $\lambda = 1$ and $\lambda = 10$

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Non-preemptive: attacks might lead to AoI reduction

Preemptive: AoI cannot be reduced

Other queueing systems:

- M/M/1 queue
- M/M/1/K queue (preemptive and non-preemptive)
- ...

Variations of attacks: only some packets are discarded (not all)

1. Kaul, S., M. Gruteser, V. Rai, and J. Kenney. "Minimizing age of information in vehicular networks". In 8th Annual IEEE Communications Society Conference on Sensor, Mesh and Ad Hoc Communications and Networks (SECON). 350-358. 2011.
2. Kaul, S. K., R. D. Yates, and M. Gruteser. "Status updates through queues". In 46th Annual Conference on Information Sciences and Systems (CISS). 1-6. 2012.
3. R. D. Yates, Y. Sun, D. R. Brown, S. K. Kaul, E. Modiano, and S. Ulukus. Age of information: An introduction and survey. IEEE Journal on Selected Areas in Communications, 39(5):1183-1210. 2021.

That's all folks

THANKS FOR YOUR ATTENTION

QUESTIONS?